Analyzing the Performance of CS 430 Project

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Abstract

Quicksort is one of the most widely used sorting algorithms. QuickSort is a Divide and Conquer algorithm. It picks an element as pivot and partitions the given array around the picked pivot. There are many different versions of quickSort that pick pivot in different ways.One is pick a random element as pivot.the other is pick median as pivot.

Median is an algorithm for selecting the kth largest element in an unordered list, having worst case linear time complexity [1]. It finds the approximate median in linear time which is then used as pivot in the quickselect algorithm. This approximate median can be used as pivot in Quicksort, giving an optimal sorting algorithm that has worstcase complexity O(n log n) [2].

Graphs were also plotted for the usual Quicksort function and the Quicksort function that uses the proposed median as pivot for relatively large values of size of array and results were compared. These confirm that proposed algorithm indeed has worst-case complexity O(n log n) for Quicksort.

**keywords**：Median , Quicksort, Partition, Median Selection

Introduction

We know select one number as the pivot from unordered list , If unordered list is partially ordered, the fixed selection of pivots makes quicksort less efficient. To alleviate this situation, pick a random element as pivot is a good idea.

Median is the middle value in a data set. Median selection is a problem that can be considered a special case of selecting the ith smallest element in an ordered set of n elements, when . An approach to solve this problem could be to sort the list and then choose the ith element. This could be using any sorting algorithm such as - Heapsort that has the worst case upper bound as O(n log n), Quicksort that has an expected running time O(n log n) though its running time is O(n2 ) in the worst case. Once the data values are sorted, it takes O(1) time to find the ith order statistics. Using an optimal sorting algorithm, the aforesaid approach gives complexity of O(n log n) as upper bound for selecting ithorder statistics [2].,A better method is pick median element as pivot .

Methodology

This report is to answer those questions:

1、In the same amount of data, the run time of quick sort with median as pivot is better than quick sort which chooses a random element in the array as the pivot

2、 Both of The time complexity of quick sort with median as pivot and The time complexity of quick sort which chooses a random element in the array as the pivot is O(nlgn）

3、In the same amount of data, the run time of order statistics median finding algorithm that use randomized is better than not use randomized

4、The time complexity of randomized median finding algorithm is O(n)

This report main method used:

In the same environment, different algorithms are tested with different data of different orders of magnitude and different characteristics, and the actual running time of the program is recorded

The main data sets include:

positive order data, reverse order data, data generated with randomized, data with a lot of number of same elements, data of different orders of magnitude

Results and Discussions

| **Median Finding** | **randomized(ms)** | **right(ms)** |
| --- | --- | --- |
| positive order 1k | 0.48 | 93.26 |
| reverse order 1k | 0.86 | 87.58 |
| random 1k | 0.63 | 0.35 |

| **Quick Sort** | **randomized(ms)** | **right(ms)** |
| --- | --- | --- |
| positive order 1k | 4.61 | 2.93 |
| reverse order 1k | 4.35 | 2.69 |
| random 1k | 4.43 | 3.15 |

| **Quick Sort** | **randomized(ms)** | **right(ms)** |
| --- | --- | --- |
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| positive order 1k | 4.61 | 2.93 |
| reverse order 1k | 4.35 | 2.69 |
| random 1k | 4.43 | 3.15 |

| **Quick Sort** | **randomized(ms)** | **Median(ms)** |
| --- | --- | --- |
| positive order 1k | 6.56 | 24.12 |
| reverse order 1k | 7.73 | 25.17 |
| random 1k | 6.73 | 26.06 |

Conclusion

Reference

1. Manuel Blum, Robert W. Floyd, Vaughan Pratt, Ronald L. Rivest, and Robert E. Tarjan. Time bounds for selection. Journal of Computer and System Sciences, 7:448{461, 1973.

[2] Introduction to Algorithms, Second Edition, Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest, Clifford Stein, The MIT Press, ISBN 0- 07-013151-1 (McGraw-Hill)